Scheduling Tasks with Precedences on Edge-Cloud Platforms Partially Powered with Renewable Energy

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General Model: Platform



Network: a complete graph with latency value for each edge

General Model: Platform



Network: a complete graph with latency value for each edge For each computing site: #CPUs #memory units speed (work/second) Pstatic Pdynamic (per CPU)

Renewable sites: prediction of power production assumed to be constant Jobs are submitted over time, their execution must start **immediately**. Tasks are allocated on **reservations** that can be changed until they start. Executing tasks cannot be preempted/migrated.



Additional Constraints: Labels

Labels are associated to tasks and computing sites for **affinity** filtering (à *la Kubernetes*).

 \rightarrow A task can only be placed on a computing site containing its labels.



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A successor task **must start its execution immediately** after all communications have finished.



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 \rightarrow A task may start **before** its actual execution (holding the resources).



Green Scheduling at the Edge

Problem recap'

- Computing sites composed of one server with limited resources
- Jobs with submission times, deadlines, DAG of tasks
- Renewable energy sources + the regular grid (brown energy)
- Labels on tasks and servers for filtering (à la Kubernetes)
- Dynamic setting with a time window of 15 minutes

Decisions to take

- Determine the allocation of tasks for each job (or reject the whole job)
- Determine the repartition of renewable/brown power drawn by each computing site for each time interval
- The tasks start/finish times are fixed by the allocation decisions

Objective: Minimize the total brown energy consumption **and** the number of jobs rejected

Some Formulae

Processing time of a task:

$$\mathsf{exec} = \frac{\mathsf{task} \; \mathsf{work}}{\mathsf{server} \; \mathsf{speed}}$$

Communication time between two tasks executed on sites i and m:

$$comm = \begin{cases} latency_{im} & \text{if } i \neq m \\ 0 & \text{if } i = m \end{cases}$$

(We assume infinite bandwidth to avoid network sharing problems)

Power consumption of a computing site when u CPUs are used:

$$power = \begin{cases} P_{static} + u \times P_{dynamic} & \text{if } u > 0 \\ 0 & \text{if } u = 0 \end{cases}$$

CPU used

















Example: Energy Consumption



Example: Energy Consumption



Example: Energy Consumption



So how to actually solve the problem?

First try with Constraint Programming

Second try using 'Classical' Scheduling Algorithms

MiniZinc

An open-source constraint modelling language which can be used in conjunction with multiple back-end solvers (*Gecode*, *OR-Tools*, etc.).

Algorithm: Upon submission of a job, instanciate a MiniZinc model and solve it to schedule all job tasks, given the current state of the platform.

Current MiniZinc models:

- Exact model: Optimal objective value (additionnal brown energy consumed) but slow to solve (several minutes)
- Approximate model: Each task consumes a fraction of the P_{static}
 Approximate objective value but fast to solve (less than a second)

Develop greedy algorithms and their variants:

- Greedy algorithms
- Local search and exhaustive search variants
- (Metaheuristics)

Rank (à la HEFT)

The rank of a task denotes its average critical path to the last finishing task of the job. For a task k:

$$\operatorname{rank}_{k} = \overline{\operatorname{exec}_{k}} + \max_{p \in \operatorname{successors of } k} \{\overline{\operatorname{comm}_{kp}} + \operatorname{rank}_{p}\}$$

 \rightarrow Used to order tasks

Candidate Locations

Define the possible placements of a task

- the allocation (the 'where')
- the anticipated start time, real start time and finish time (the 'when')
- some metrics (additional total/brown/renewable energy consumption)

The locations are *feasible*: they respect the resource demands and deadline

Example: Candidate Location



Power used (W)



Example: Candidate Location



 Candidate 1: Cloud site Start time: 5, finish time: 5.5 70J additional brown energy (P_{static} + P_{dynamic})

Example: Candidate Location



- Candidate 1: Cloud site Start time: 5, finish time: 5.5 70J additional brown energy (P_{static} + P_{dynamic})
- Candidate 2: Edge site
 Start time: 6, finish time: 7
 25J additional brown energy (P_{dynamic} only)

Algorithm: Greedy

Greedy Algorithm Skeleton

Upon job submission:

- Sort its tasks by **decreasing rank** (follows the dependencies)
- 2 For each task:
- 3 Compute all feasible candidate locations
- 4 Place the task on the **best location**

If a task cannot be allocated (lack of free resource or deadline not met) \rightarrow Reject the whole job

Selection policy for the best location:

- Smallest additionnal energy
- Smallest additionnal brown energy
- Some other combination (resource usage, energy, finish time)

Example: Greedy Algorithm



Example: Greedy Algorithm



Example: Greedy Algorithm

- \rightarrow New job submission at time 3 One task, Edge only, 3 CPUs
- No space for T_3
 - Reject the job (Greedy algorithm)
 - Try to rearrange the schedule (Local Search)









Local Search procedure

- Make a backup of the platform state
- Remove all non-running tasks in the "neighborhood" of the problematic task
- **3** Schedule the problematic task
- (Try to) schedule all tasks in the queue using a scheduling procedure



Scheduling procedures

- **Greedy**: Simply apply the greedy algorithm.
- Greedy with LS: Perform greedy but apply the Local Search procedure if necessary.
- Exhaustive search: Recursively try all candidate locations for all tasks. Apply the best solution found.

 \rightarrow If a task cannot be allocated, revert to backed up state and reject the job of first problematic task.

Not really started yet

(but algorithms are implemented)

- Try different types of platforms
- Try different collections of jobs

For that we need datasets (maybe look at *Azure* ones?)

 \rightarrow Do you have ideas of datasets? Please come and talk to me



A classical scheduling problem with uncommon constraints:

- Renewable power sources
- Allocation filtering with labels
- Immediate starting times (no delay)

On-going and future work:

- Focus on the experimental part
- Continue the algorithm design/implementation with new ideas
- Design good search heuristics for the Constraint Programming models

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Site 4











